EECS 388: Computer Systems and Assembly Language

Homework 2 - Solution

Justify your answers!

Figure 1 shows a part of the memory (both contents and locations).

Contents L	Contents Locations					
\$20	\$4000					
\$50	\$4001					
\$01	\$4002					
:	:					
:	:					
\$B5	\$5000					
\$CD	\$5001					
Figu	re 1.					

Problem 1 (10 Points):

Consider the memory shown in Figure 1. What is in accumulator A and the N, Z, C, V bits in CCR after an LDD #\$4001 instruction is executed?

Accumulator A _____\$40___ N __0_ Z __0_ C __?(unchanged)_ V __0__

LDD #\$4001 →

This instruction will load register D (A:B) with the immediate value of \$4001

D = (A:B) = \$4001, therefore... A = \$40 B = \$01

According to pg. 450 of the textbook the LDD instruction has an effect on the N and Z bits, it will zero out the V bit, and it has no effect on the C bit.

The N-bit is set when a number is negative in 2's complement representation. This means that this bit will be set if the most significant bit of D is set. The most significant nybble (4-bits) of D is 4, or 0100 in binary, so the most significant bit is 0, and thus the N-bit is also zero.

The Z-bit is set when the result of an operation is zero. In this case, Z will be zero due to the fact that the number being loaded into register D is not a zero.

Problem 2 (15 Points):

Write a program segment to reverse the bit order of a 6-bit number. Assume this number is stored in \$6000. Store the reversed number to \$6001 (i.e., if the original number is 0 0 b₅ b₄ b₃ b₂ b₁ b₀, after this program, the number in \$6001 will be 0 0 b₀ b₁ b₂ b₃ b₄ b₅).

Reversing a number can be easily accomplished by rotating the number 1-bit at a time, until the desired rotation is accomplished. First, start the process by rotating the LSB out of the original number and into the carry register. Then rotate this number into the LSB of another register, and repeat this process. Each subsequent rotation will "push" the last rotation forward until all bits are in their correct places. This can be accomplished with a series of shifts, or with a loop.

ORG \$4000	; Set program start address			
LDAA \$6000 LDAB #\$00	; Load original number in register A ; Initialize register B			
LSRA ROLB				
; Now regis	ster B = $000000b_0$			
LSRA ROLB	; Shift LSB out of original number and into carry register ; Rotate bit in carry register into LSB of register B			
; Now register B = 000000b₀b₁				
LSRA ROLB	; Shift LSB out of original number and into carry register ; Rotate bit in carry register into LSB of register B			
; Now regis	ster B = $00000b_0b_1b_2$			
LSRA ROLB	; Shift LSB out of original number and into carry register ; Rotate bit in carry register into LSB of register B			
; Now regis	ster B = $0000b_0b_1b_2b_3$			
LSRA ROLB	; Shift LSB out of original number and into carry register ; Rotate bit in carry register into LSB of register B			
; Now regis	ster B = $000b_0b_1b_2b_3b_4$			
LSRA ROLB	; Shift LSB out of original number and into carry register ; Rotate bit in carry register into LSB of register B			
; Now register $B = 00b_0b_1b_2b_3b_4b_5$				
STAB \$6001 SWI	; Store the results and stop the program			

Problem 3 (20 points):

Write a program segment to multiply a 16-bit number in the D register by 10 using arithmetic left shift instead of multiplication instructions.

This problem can be solved using simple algebra techniques. Let the 16-bit number be called X:

 $10^{*}X$ = (8 + 2)*X = (8*X) + (2*X) = (arithmetic shift left X 3 times) + (arithmetic shift left X once)

; Set program start address
is already loaded into D
ne shift) Save 2*X
vo more shifts shift)
000 contains 2*X and register D holds 8*X I the two numbers together
) = D + MEM[\$6000]
program

Problem 4 (20 points):

Write a program to subtract two 24-bit numbers and store the result to memory locations starting at \$6000. The two 24-bit numbers are stored in memory locations starting at \$5000 and \$5010, respectively.

The process of subtraction requires one to perform carrying if needed. Keeping track of carry bits is difficult, so an easier solution is to simply negate one of the numbers and add b/c algebraic rules tell us that X + Y is the same as X + (-Y).

The process of negating a number (in 2's complement) requires one to negate all of it's bits and add one. The HC12 does not have a simple negation operator (it does have a NEG operation which converts an 8-bit number to negative 2's complement, but this can't be used for 24-bit numbers!!) so instead we can use a trick. The trick is

that a bit can be negated by using an exclusiveOR with the binary value '1'. Therefore a full 8-bit negation can be accomplished by exclusiveOR with \$FF.

Not (X) = ExclusiveOR(X,\$FF)

```
NEGATE MASK
                    EQU
                                   ; Create an equate for the negation mask
                           $FF
BASE1
                                   $5000 ; Base of 1st number
                           EQU
BASE2
                           EQU
                                   $5010 ; Base of 2nd number
NEGATED BASE
                    EQU
                           $5020 ; Base of negated version of 2nd number
RESULT_BASE
                           EQU
                                   $6000 ; Base of result
       ORG $4000
                           ; Set program start address
       ; Convert the second number to negative in 2's complement
       ; First Negate the 2nd number 1 byte at a time
       LDAA BASE2+0
                                   ; Read in byte
       EORA #NEGATE MASK+0
                                   ; Negate it
       STAA NEGATED_BASE+0
                                   ; Store the result
       LDAA BASE2+1
                                   ; Do this for the remaining 2 bytes
       EORA #NEGATE MASK
       STAA NEGATED BASE+1
       LDAA BASE2+2
       EORA #NEGATE MASK
       STAA NEGATED_BASE+2
       ; Now perform 2's complement addition
       ; NUM1 - NUM2 = NUM1 + (Negated NUM2) + 1
       ; Pre-load the carry bit with 1
       LDAA #$01
                    ; Load a 1 into A, and shift the 1 into the carry reg
       LSRA
       ; Now begin the addition starting with the least-significant byte
       LDAA BASE1+2
                                   : Load in byte of NUM1
       ADCA NEGATED_BASE+2
                                   ; Add negated byte with carry
       STAA RESULT_BASE+2
                                   ; Use carry (with carry pre-set to 1)
       LDAA BASE1+1
                                   ; Continue for remaining bytes
       ADCA NEGATED_BASE+1
       STAA RESULT BASE+1
                                   ; Use carry again
       LDAA BASE1+0
       ADCA NEGATED BASE+0
       STAA RESULT_BASE+0
                                   ; Use carry again
       SWI; Stop the program
```

Problem 5 (20 Points):

If A contains \$56, what is the result of each of the following instructions? Assume that A is restored to its original value before each instruction.

```
a) ANDA #$33
A = $56
           = 0101 0110
Mask = $33 = 0011 0011
Result(AND) = 0001 0010 = $12
A = $12
b) ORAA #$33
A = $56 = 0101_0110
Mask = $33 = 0011 0011
Result(OR) = 0111 0111 = $77
A = $77
c) EORA #$33
A = $56
           = 0101 0110
Mask = $33 = 0011 0011
Result(EOR) = 0110 0101 = $65
A = $65
```

d) BITA #\$80

The BITA operation performs arithmetic, but does not alter the contents of register A. However, it does alter the contents of the CC register. Typically this instruction will be used for the sake of comparison before a branch instruction, so that CCR bits can be purposefully set w/o having to alter any of the values in registers.

Problem 6 (15 Points):

Consider the following program:

LDD #\$F00D STD \$8100 BSET \$8100, \$44 BCLR \$8101, \$11

What numbers are in \$8100 and \$8101 at the end?

LDD #\$F00D

This instruction loads register D (A:B) with the immediate value \$F00D.

D = \$F00D \$8100= ?? \$8101= ??

STD \$8100 This instruction writes the contents of D, which are \$F00D, to memory location \$8100

D = \$F00D \$8100= \$F0 \$8101= \$0D

BSET \$8100, \$44 This instruction alters memory location \$8100 by setting bits that are 1 in the mask.

\$8100 = \$FO = 1111_0000 Mask = \$44 = 0100_0100 Result = 1111_0100 = \$F4

D = \$F00D \$8100= \$F4 \$8101= \$0D

BCLR \$8101, \$11 This instruction alters memory location \$8101 by clearing bits that are 1 in the mask.

\$8101 = \$0D = 0000_1101 Mask = \$11 = 0001_0001 Result = 0000_1100 = \$0C

D	= \$F00D			
\$8100)= \$F4			
\$810 1	1 = \$0C			